

# Multicoloured cliques in vertex-coloured graphs

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A **vertex-colouring** of graph  $G$  is a function  $c: V(G) \rightarrow [k(c)]$ . A **vc-graph** is a pair  $(G, c)$ .

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$S \subseteq V(G)$  is **rainbow** for colouring  $c: V(G) \rightarrow [k(c)]$  if  $|c(S)| = |S|$ , **strong rainbow** if further  $|S| = k(c)$ .

IS

Input: graph  $G$ , integer  $k$

Question: does  $G$  contain independent set  
of size  $k$ ?

$IS(\mathcal{C})$

Input: graph  $G$ , integer  $k$ ,  $G \in \mathcal{C}$

Question: does  $G$  contain independent set  
of size  $k$ ?

ISGCI: 1031 classes at present

**Theorem** (Chen, Huang, Kanj, Xia, 2006): if IS can be solved in  $f(k)m^{o(k)}$  time then it can be solved in  $2^{o(n)}$  time.

## SRIS

Input:  $\text{vc-graph } (G, c)$

Question: does  $(G, c)$  contain a strong rainbow independent set?

## SRIS

Input:  $\text{vc-graph } (G, c)$

Question: does  $(G, c)$  contain a rainbow independent set of size  $|c(V(G))|$ ?

## RIS

Input: vc-graph  $(G, c)$ , integer  $k$

Question: does  $(G, c)$  contain a rainbow independent set of size  $k$ ?

# Multicoloured cliques in vertex-coloured graphs

# Multicoloured **independent sets** in vertex-coloured graphs

# Strong rainbow independent sets in vertex-coloured graphs

Suppose  $S = (V, (Q_i)_I)$ ,  $T = (W, (R_i)_I)$  are relational structures with  $r(Q_i)_I = r(R_i)_I$ .

$f$  is **homomorphism** from  $S$  to  $T$  if  $f: V \rightarrow W$  such that  $(u_1, u_2, \dots, u_{r(Q_i)}) \in Q_i \Rightarrow (f(u_1), f(u_2), \dots, f(u_{r(Q_i)})) \in R_i$ .

$S \rightarrow T \equiv \exists$  homomorphism from  $S$  to  $T$

CSP<sub>2</sub>

Input:           arity  $\leq 2$  relational structures  
                   $(S, T)$

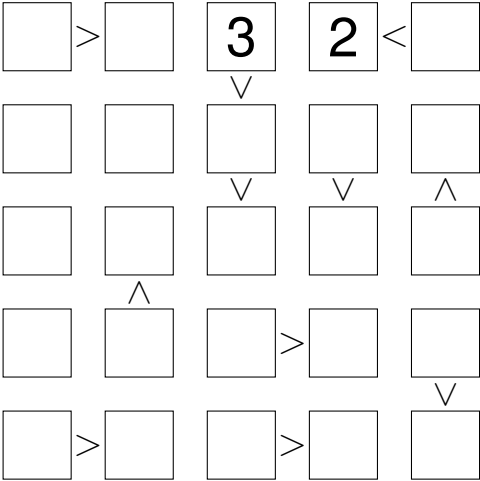
Question:     $S \rightarrow T$  ?

$\text{CSP}_2(\mathcal{C}, \mathcal{D})$

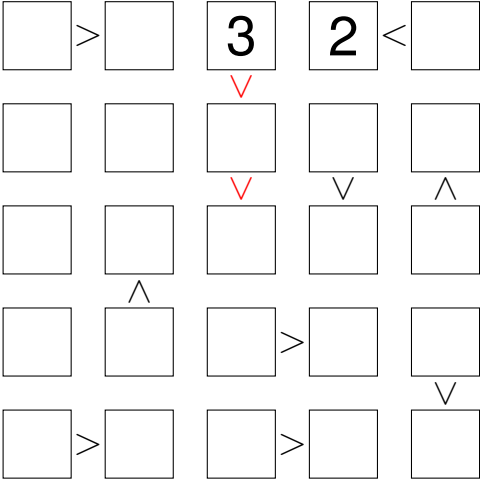
Input:           arity  $\leq 2$  relational structures  
                   $(S, T), S \in \mathcal{C}, T \in \mathcal{D}$

Question:        $S \rightarrow T$  ?

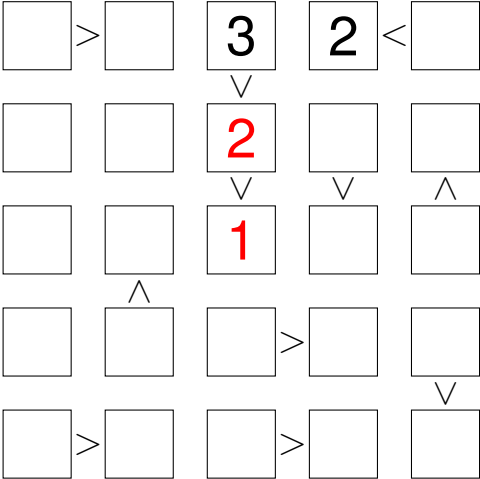
# Futoshiki



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$$([d]^2, N, L, U_{1,3}, U_{1,4}) \rightarrow ([d], \neq_d, <_d, \{(3)\}, \{(2)\})$$

$$N = \{((i, j), (i, b)) \mid j \neq b\} \\ \cup \{((i, j), (a, j)) \mid i \neq a\}$$

$$L = \{((i, j), (a, b)) \mid (i, j) < (a, b)\}$$

$$U_{i,j} = \{((i, j))\}$$



**Lemma** (Fellows, Hermelin, Rosamond, Vialette, 2008):  $k$ MC is  $W[1]$ -complete.

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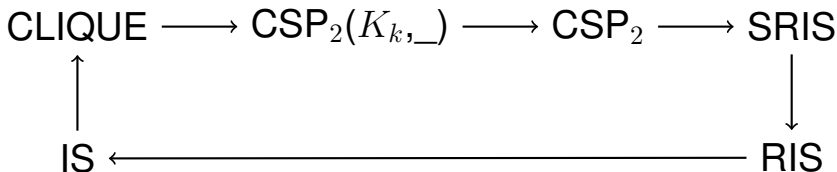
CLIQUE

$k$ MC

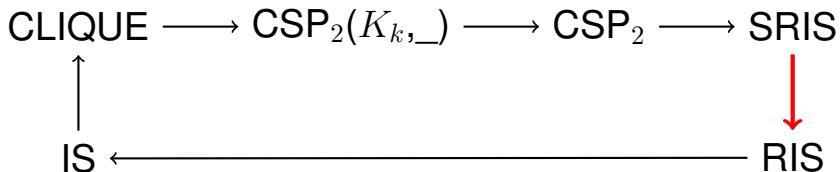


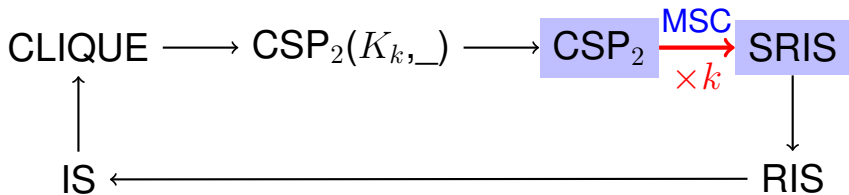
**Lemma** (Fellows, Hermelin, Rosamond, Vialette, 2008):  $k$ MC is  $W[1]$ -complete.

**Corollary:** SRIS is  $W[1]$ -complete.



Hell 1972 (graphs), Gregory Gutin and colleagues 2006 (digraphs)



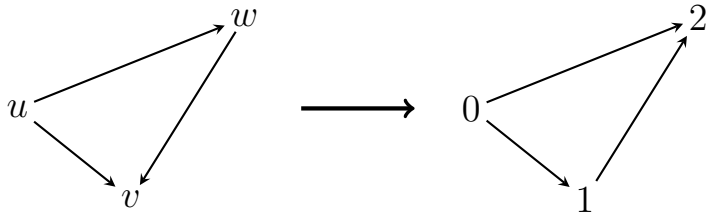


## microstructure complement

$$S = (\{u, v, w\}, \{(u, v), (u, w), (w, v)\})$$

$$T = (\{0, 1, 2\}, <_3)$$

$$S \rightarrow T ?$$

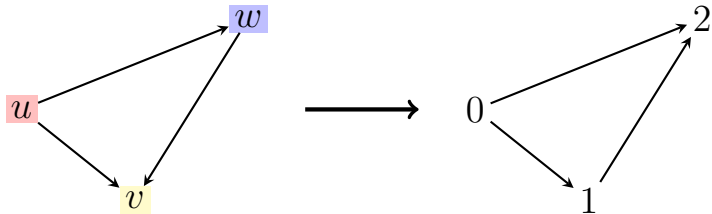


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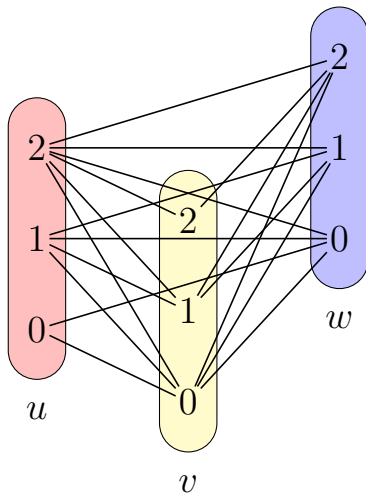
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$$S \rightarrow T ?$$



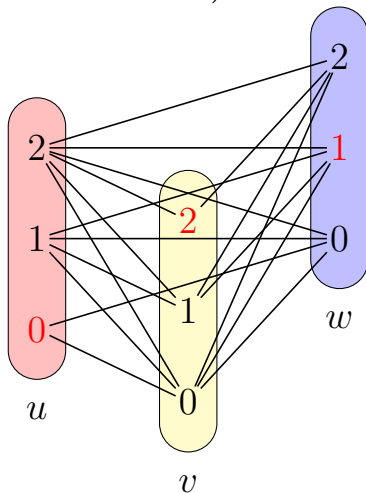
# MSC

$$(G(S \times \overline{T}), \pi_1)$$

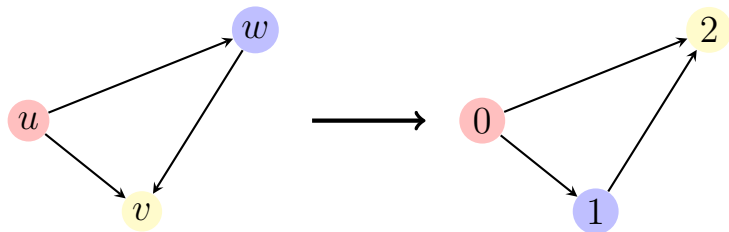


# MSC

$$SRIS((G(S \times \bar{T}), \pi_1))$$



# MSC



Rainbow

CSP

F-free

Conclusion

# MSC



$$(V, (Q_i)_I), (W, (R_i)_I) \quad (S, T)$$

# MSC



$$(V, (Q_i)_I), (W, (R_i)_I) \quad (S, T)$$

$$(V \times W, (Q_i \times \overline{R_i})_I) \quad S \times \overline{T}$$

# MSC

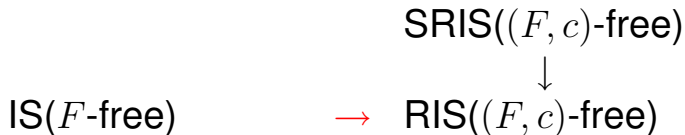


$$(V, (Q_i)_I), (W, (R_i)_I) \quad (S, T)$$

$$(V \times W, (Q_i \times \overline{R_i})_I) \quad S \times \overline{T}$$

$$(V \times W, E(\bigcup_{i \in I} Q_i \times \overline{R_i})) \quad G(S \times T)$$

dragons: need to keep  $\overline{R}$  small; universe must be finite



$F$ -free: no  $F$  induced substructure

$(F, c)$ -free: no  $(F, c)$  induced substructure

$\forall c$

SRIS( $(F, c)$ -free) NPC  
 $\Downarrow$   
IS( $F$ -free) NPC  $\Rightarrow$  RIS( $(F, c)$ -free) NPC

$F$ -free: no  $F$  induced substructure

$(F, c)$ -free: no  $(F, c)$  induced substructure

$\exists c$

$$\text{IS}(F\text{-free}) \text{ PTIME} \leftarrow \text{RIS}((F, c)\text{-free}) \text{ PTIME} \xrightarrow{\uparrow} \text{SRIS}((F, c)\text{-free}) \text{ PTIME}$$

$F$ -free: no  $F$  induced substructure

$(F, c)$ -free: no  $(F, c)$  induced substructure




**Proposition** (S. and Jeavons, 2008):




$(S, T) \in \text{ALLDIFF} \Rightarrow \text{MSC}(S, T) \in \text{PERFECT}$

**Proposition** (S. and Jeavons, 2008):

$(S, T) \in \text{TREE} \Rightarrow \text{MSC}(S, T) \in \text{PERFECT}$



IS(  -free) NPC

$\Rightarrow$  RIS( (  ,  ,  )-free) NPC

$\Rightarrow$   $\left\{ \begin{array}{l} \text{RIS(  -free) NPC,} \\ \text{RIS(  -free) NPC,} \\ \text{RIS(  -free) NPC} \end{array} \right.$



**Theorem** (Balas and Yu, 1989): IS( $2K_2$ -free)  
PTIME


**Proposition** (Cooper, Jeavons, S., 2008):

SRIS(,  -free) PTIME

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PTIME

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SRIS(, -free) PTIME

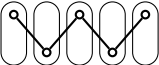
**Question:** RIS(-free) PTIME?

$\exists c$  RIS( $(P_5, c)$ -free) PTIME

$\Rightarrow$  IS( $P_5$ -free) PTIME

$\exists c \text{ RIS}((P_5, c)\text{-free}) \text{ PTIME}$

$\Rightarrow \text{IS}(P_5\text{-free}) \text{ PTIME}$

**Question:**  $\text{RIS}(\text{$

# Further work



- ▶ Detailed analysis of all reductions
  - ▶ ISGCI for vc-graphs?
  - ▶ How to use colours for kernelization?
  - ▶ Non-random colours for colour coding?
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